A Cost Model for Network Traffic (with an application to paid-peering)

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with

Murtaza Motiwala, Nick Feamster (Georgia Tech),
Anukool Lakhina (Guavus) [Cost Model]

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[Paid-peering]

Outline

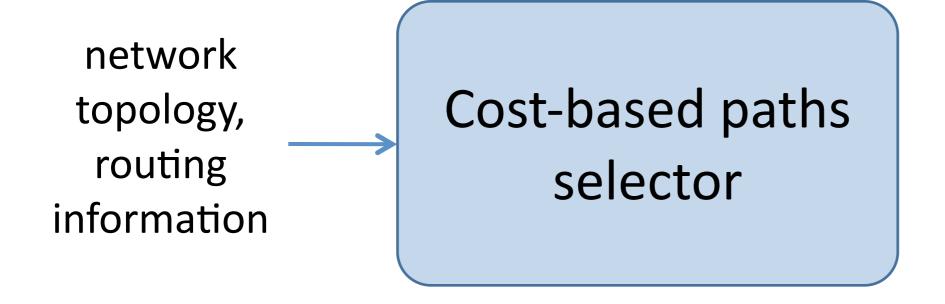
- These are really two talks
- But they are related
- Part 1: Formalizing a cost model for network traffic
 [CCR'12]
- Part 2: A value-based framework for peering agreements
 [ITC'10, NANOG 49]
- The cost model can be useful for measuring peering "value"

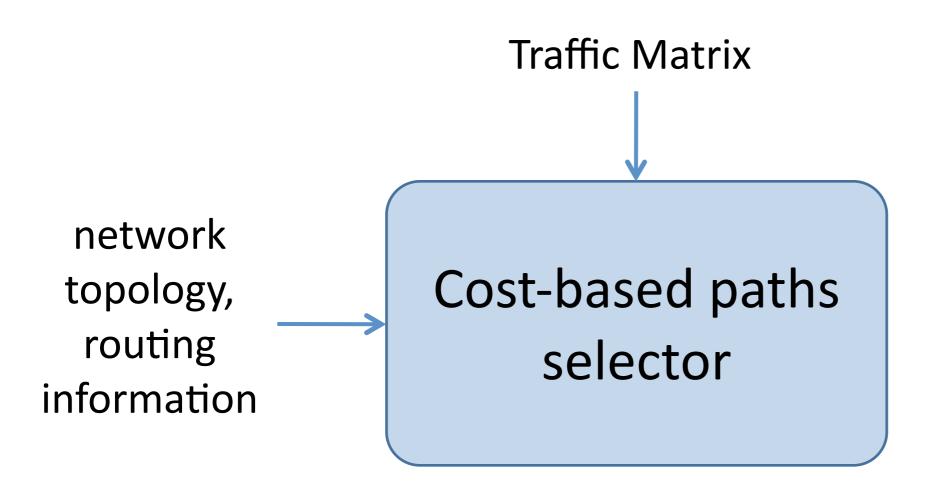
Part 1: Formalizing a Cost Model for Network Traffic

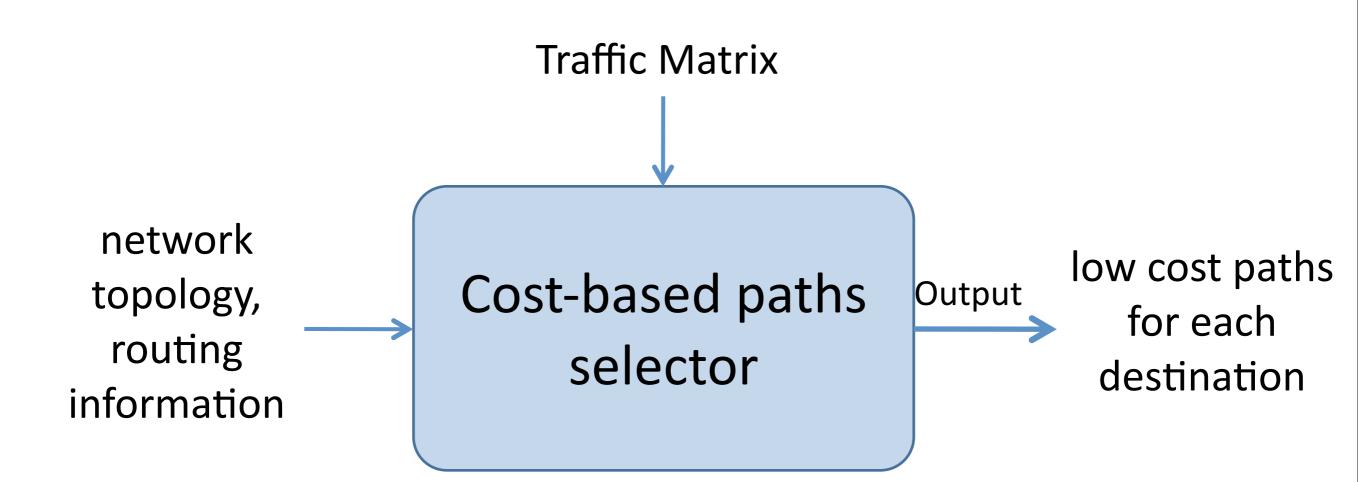
Optimizing Network Costs

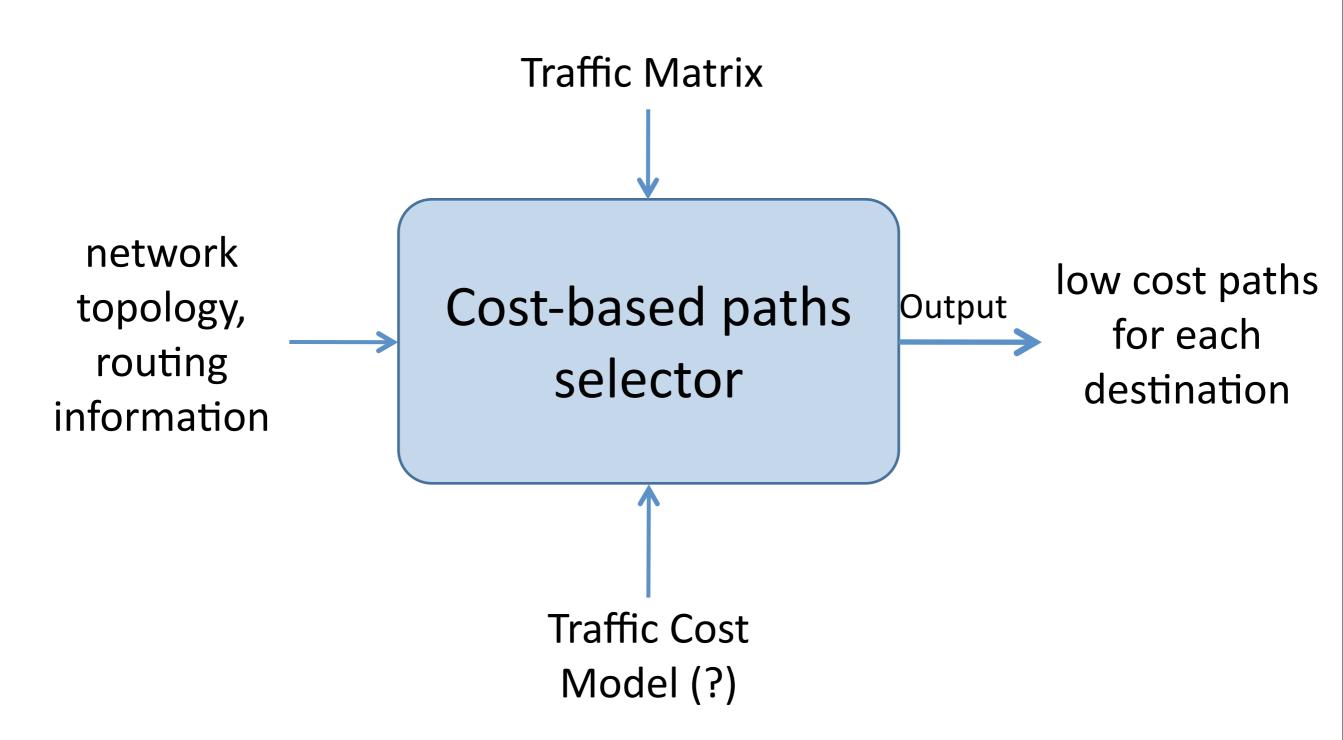
- Traffic-related costs contribute to the total cost of running a network
- Routing in recession: configuring routing in a network to minimize traffic-related costs
- Relatively easy: How do Individual elements contribute to costs? Harder: How much do individual ingress-egress flows cost?
- Need a holistic traffic cost model that can attribute costs to individual flows

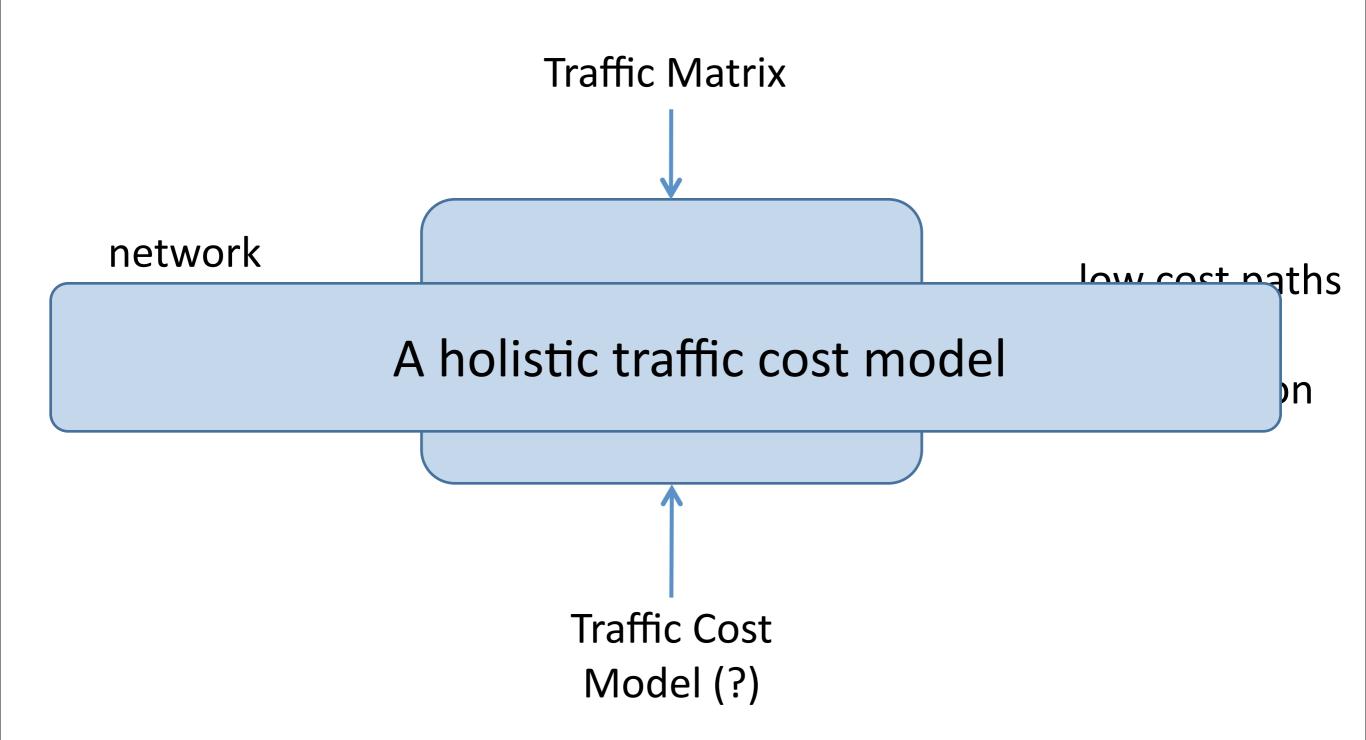
Cost-based paths selector











Costs for operating a network

- Traffic costs
 - Paying for transit, port costs, cost for laying fiber
- Operational costs
 - Paying salaries to employees
- Equipment and maintainenance costs
 - Buying networking gear, service fees to vendors
- Miscellaneous costs
 - IT related, real-estate, etc

Costs for operating a network

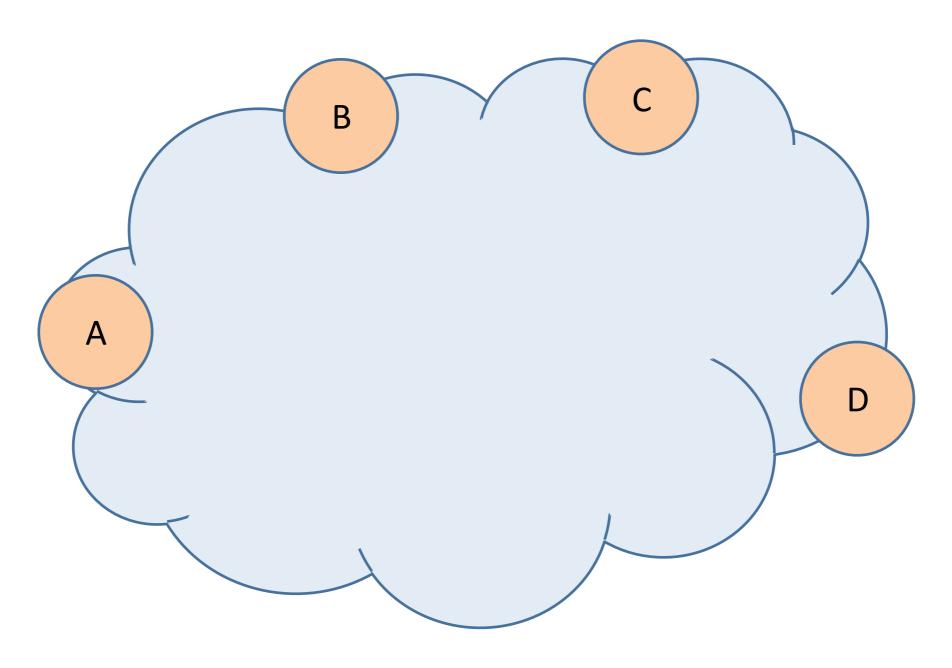
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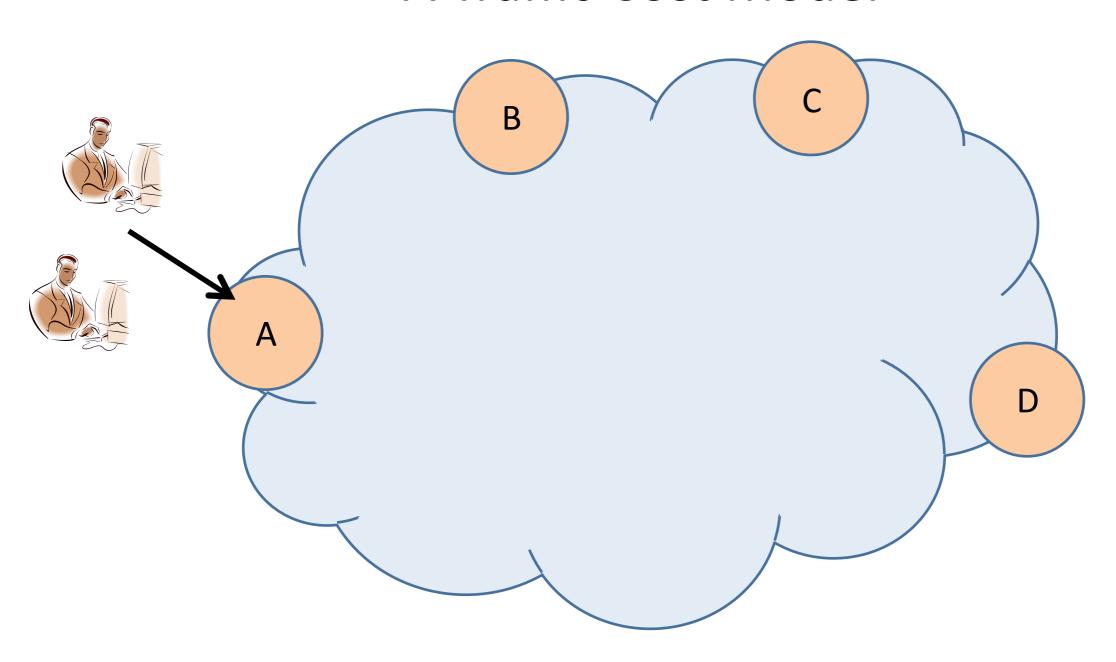
Goal: A simple but still useful cost model

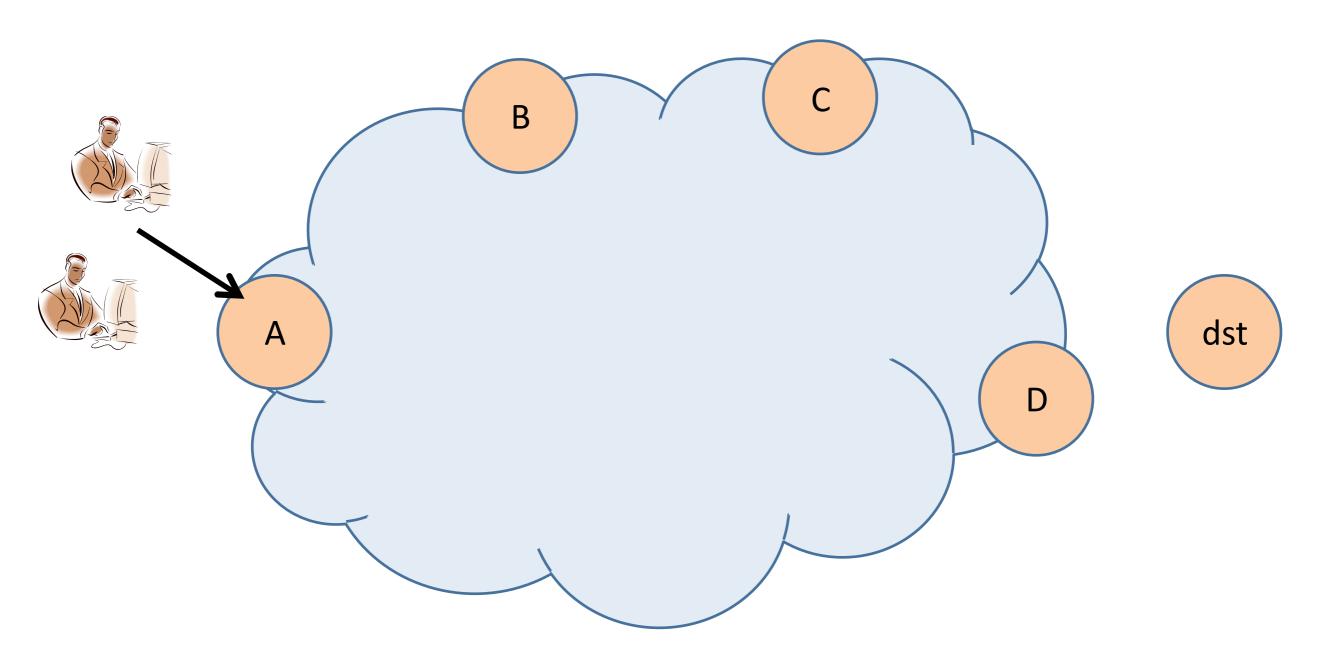
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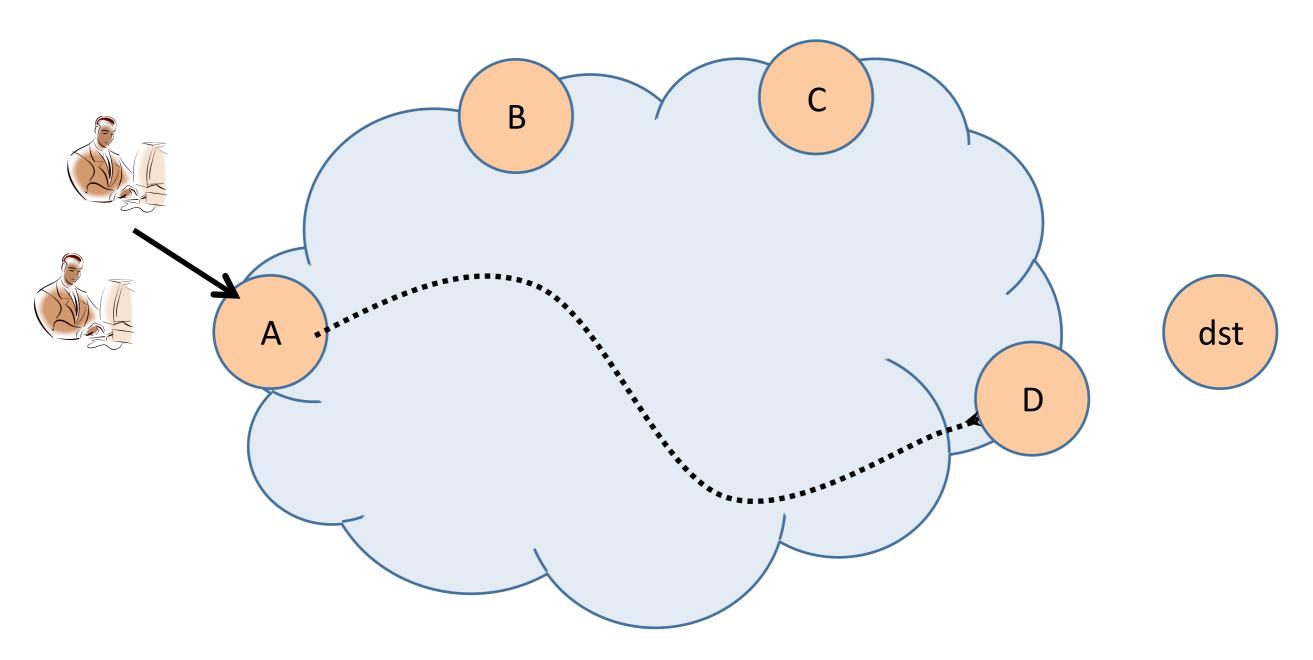
Applications

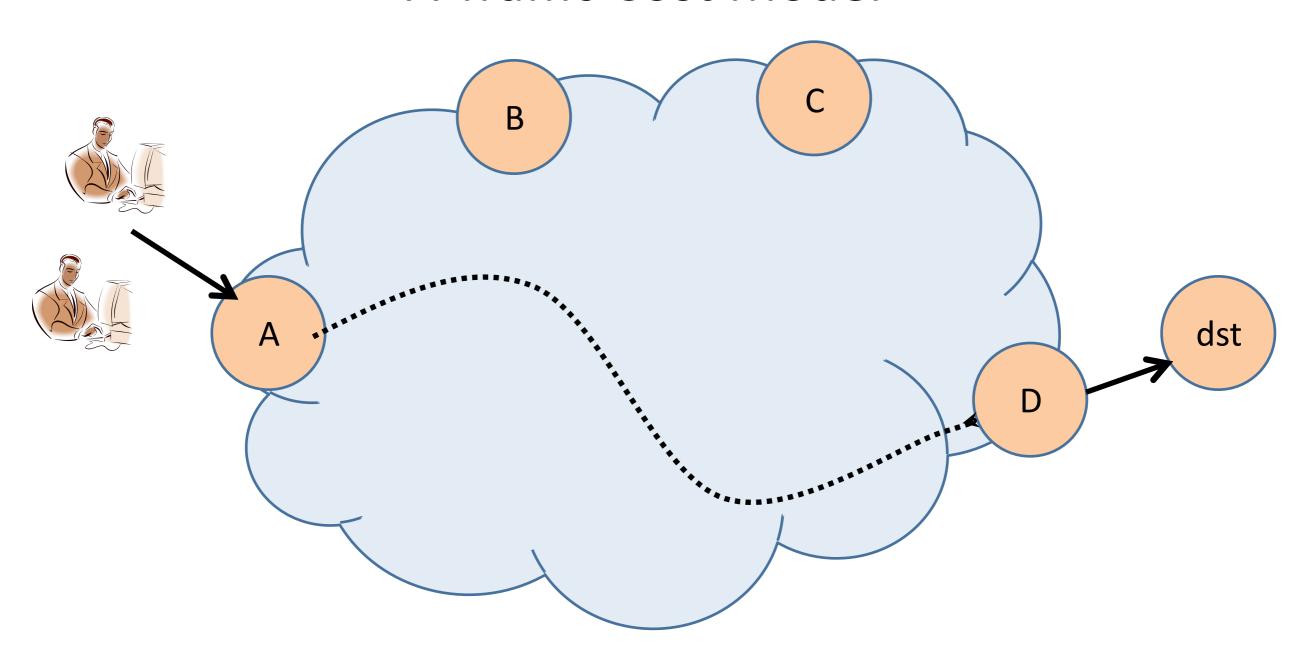
- Min-cost routing: Optimal routing of ingress-egress flows to minimize cost
- Peering Location selection: Which location to establish peering with a neighbor?
- Peering evaluation: What is the "value" of a peering link?

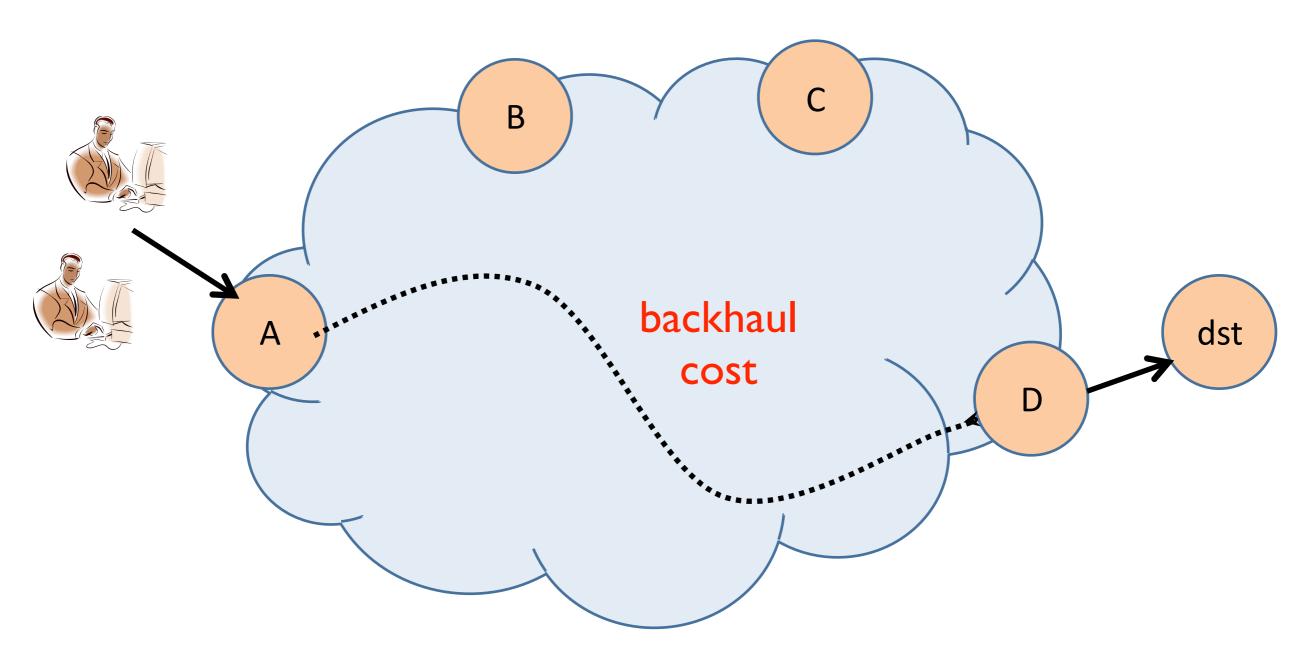


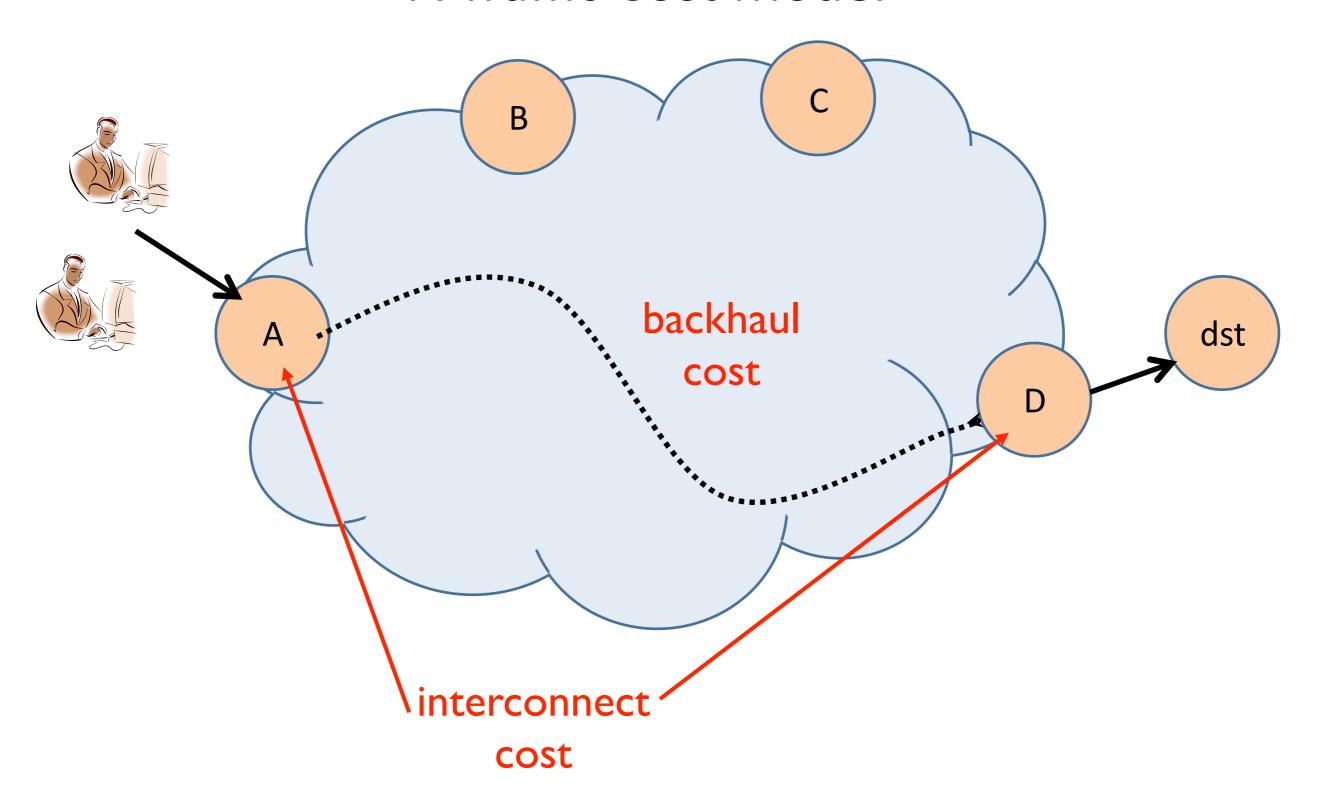












$$C_N = C_F + C_U$$

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Fixed Backhaul cost for all (PoP, PoP) pairs

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p1,p2
a,p

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Ingress Interconnect Cost

a,p

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$$C_U = \sum_f (c_i(f) + c_b(f) +$$

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Backhaul Cost

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a,p

Fixed Backhaul cost for all (PoP, PoP) pairs Fixed Interconnect cost for all (AS, PoP) pairs

Ingress-egress flows

$$C_{U} = \sum (c_{i}(f) + c_{b}(f) + c_{e}(f))$$

$$c_{\mathsf{b}}(\mathsf{f})$$

$$c_e(f)$$

Ingress Interconnect Cost

Backhaul Cost

Egress Interconnect Cost

Cost Optimization

- We focus on optimizing traffic-dependent costs
- Requires an operator to determine the cost associated with each ingress-egress flow
- Interconnect costs based on 95th percentile of total volume at that interconnect: cost = per_bit_price * V_95
- Approach 1: Assume V_95 is linear function of average rate
 - Flow's contribution = per_bit_price * constant * flow_rate
- Approach 2: Use Shapley Value
 - Flow's contribution = Shapley value across all flows at that interconnect

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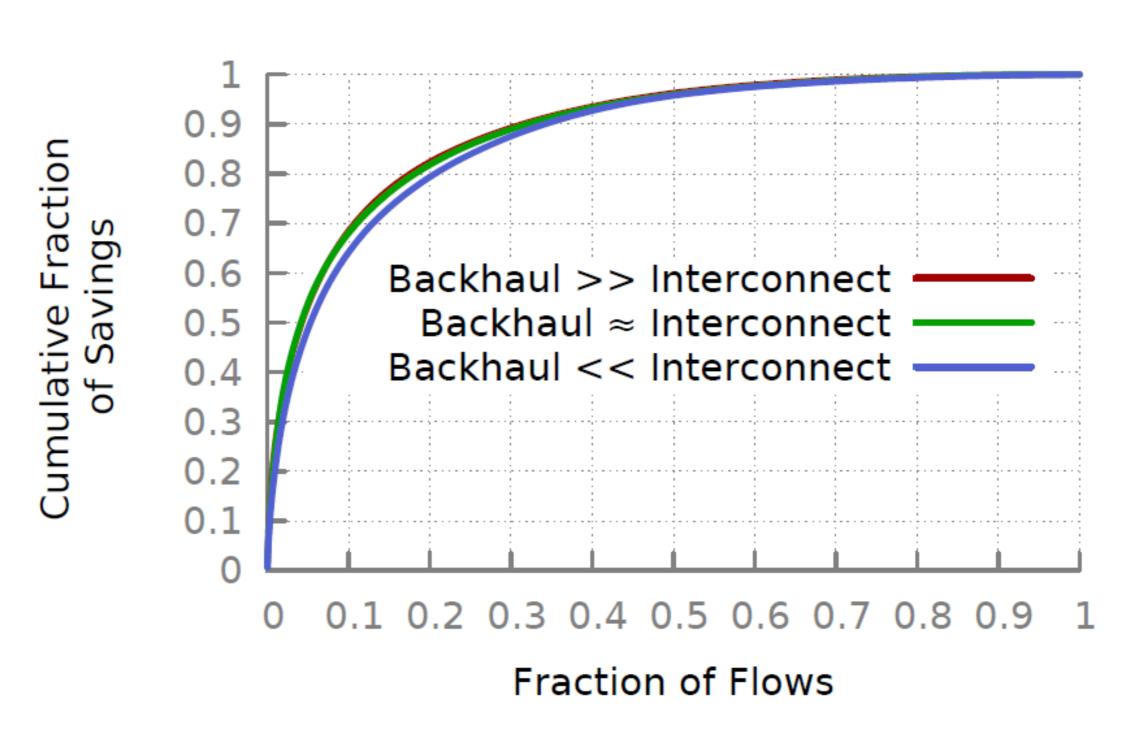
Computationally expensive!

Flow's contribution = Shapley value across all flows at that interconnect

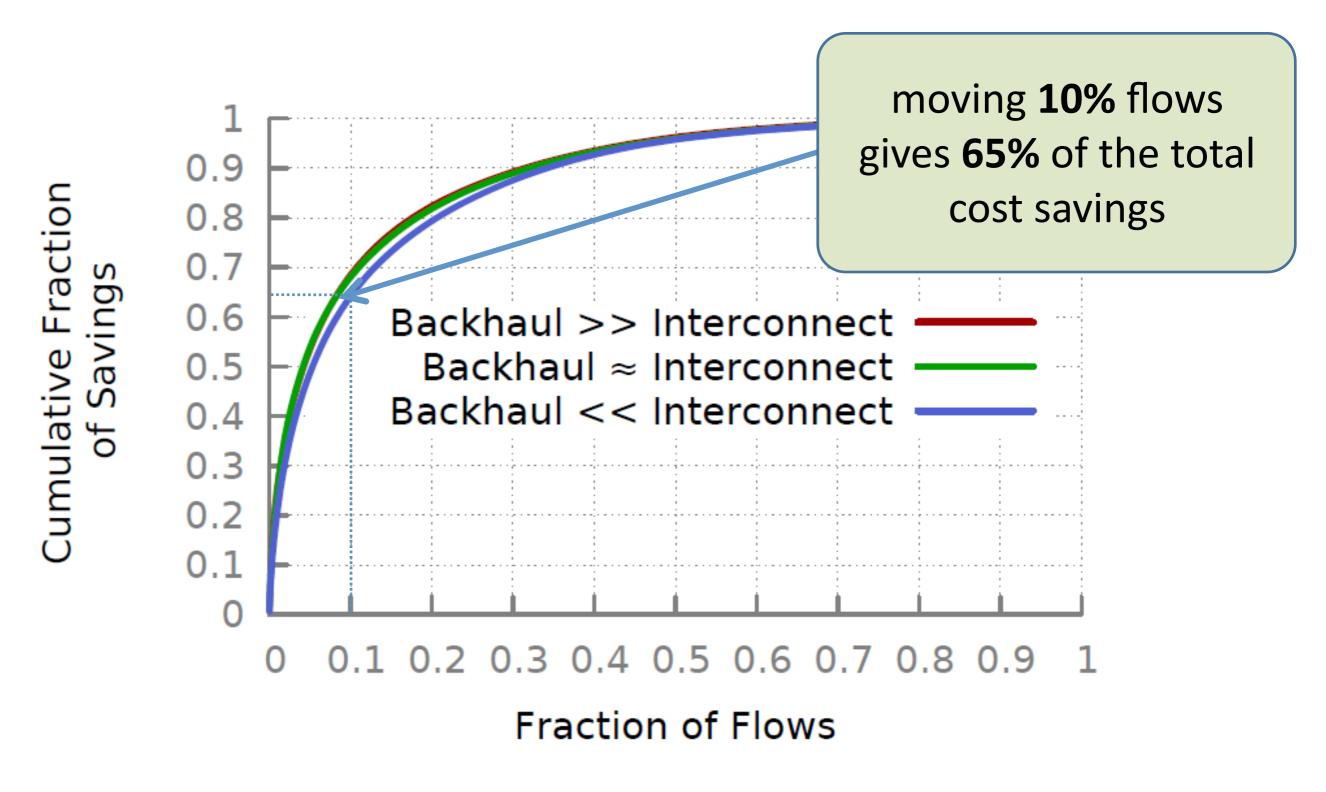
Application 1: Minimize Cost of Routing Traffic

- Objective: Minimize total cost of all ingress-egress flows
- Constraints: Backhaul, interconnect link capacities
- Knobs: egress (pop, AS) for each flow
- NP-hard to determine optimal routing! Use a greedy algorithm to approximate the optimal solution
- Iteratively, for each flow f in descending order of flow costs
 - For each PoP p, find the cheapest AS at p which has a route to f's destination
 - Assign f to the cheapest egress (PoP, AS) pair

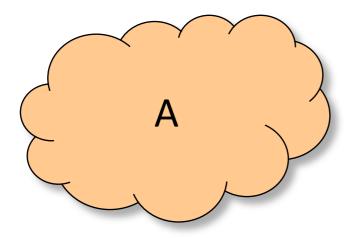
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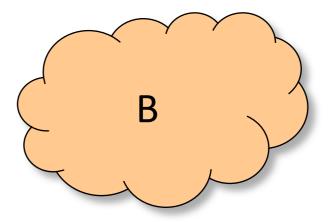


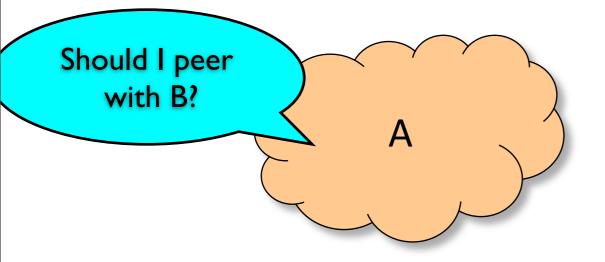
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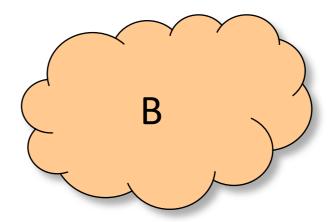


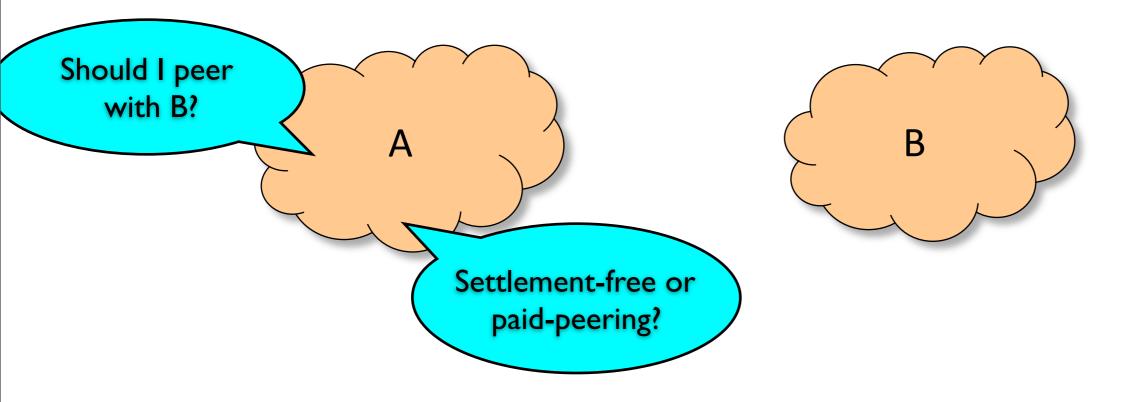
Part 2: A Value-based Framework for Peering Decisions

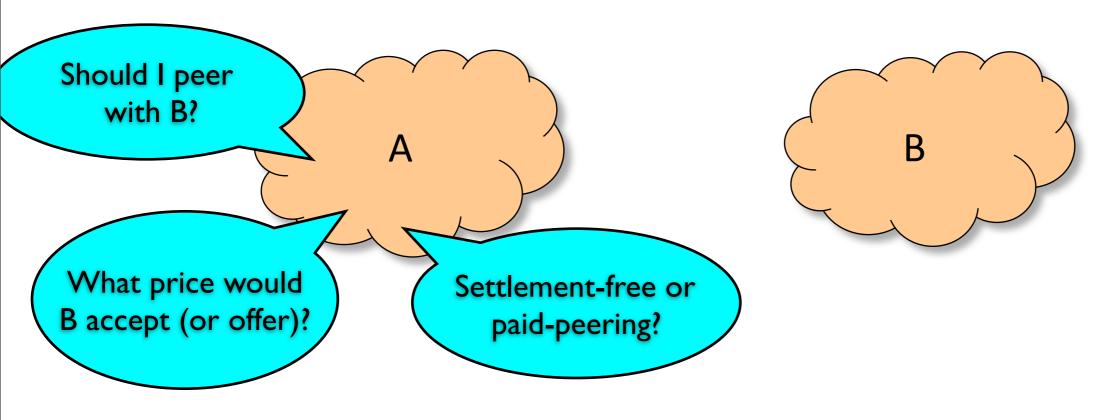


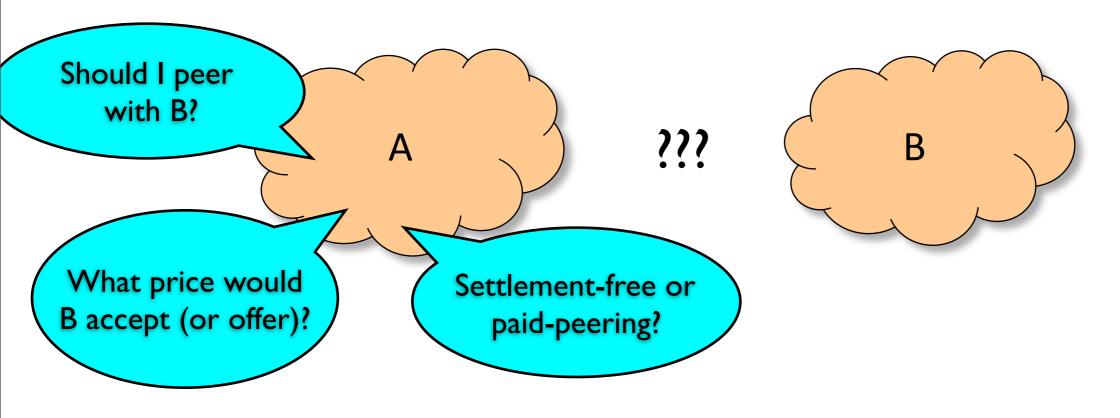






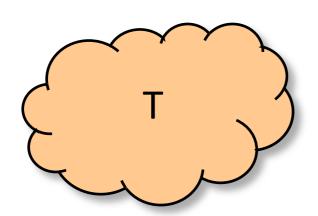


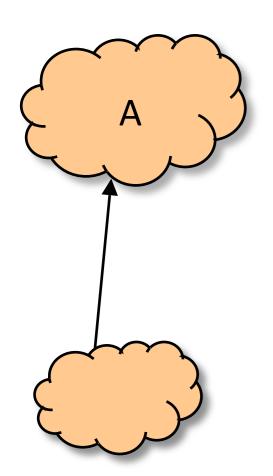


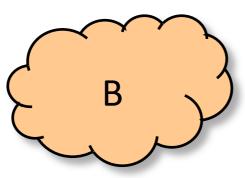


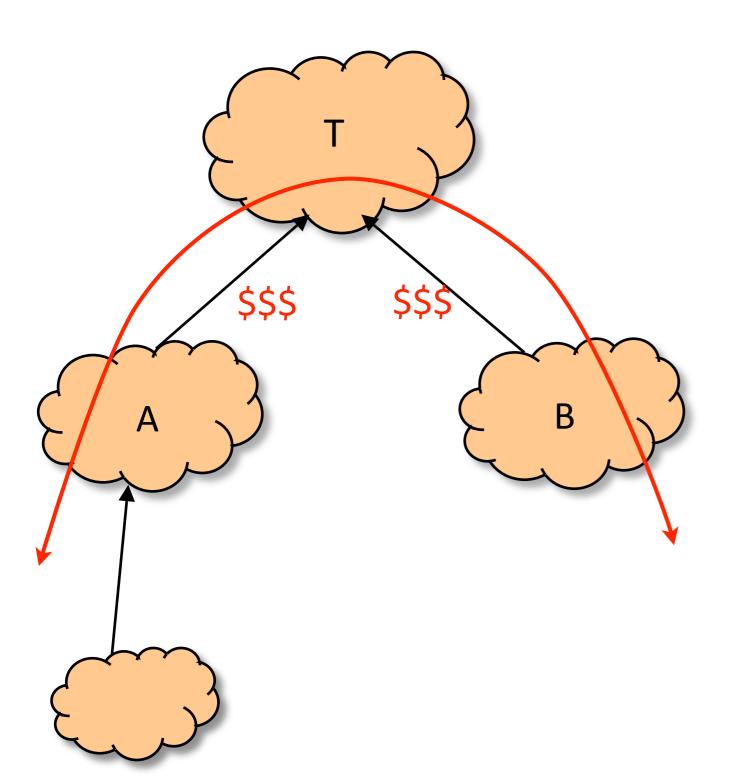
Value Based Peering

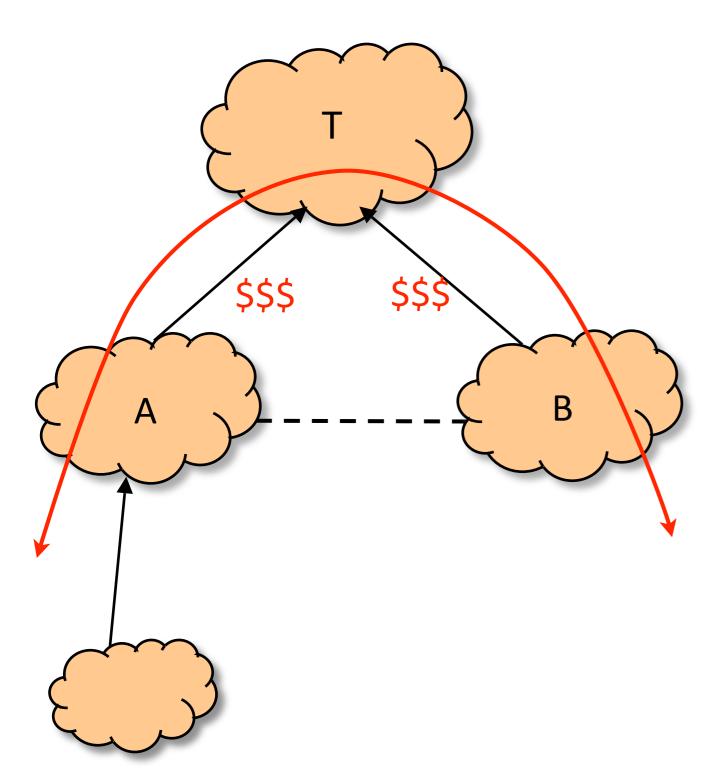
- Price based on the "value" of the link
- For a network, define the notion of economic "fitness"
- fitness = revenue interconnect costs backhaul cost
- Value of a peering link is the difference in fitness with and without the link
- Value = f_{with} $f_{without}$
- Revenue and costs could change on peering/depeering

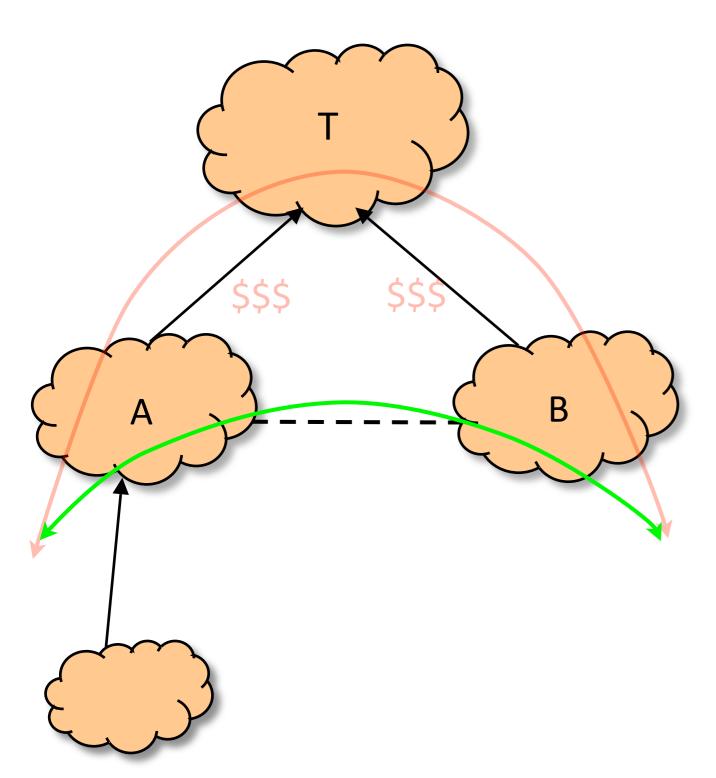




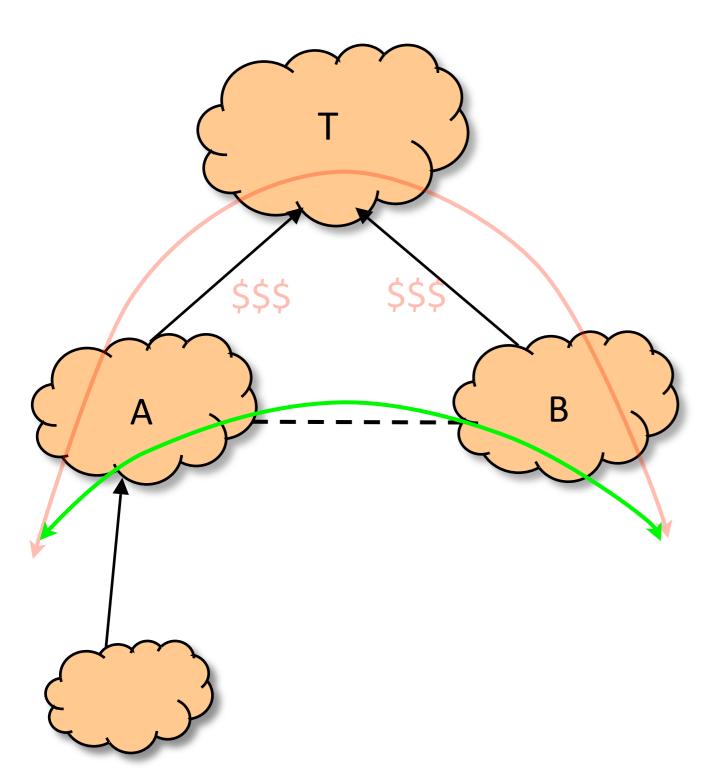




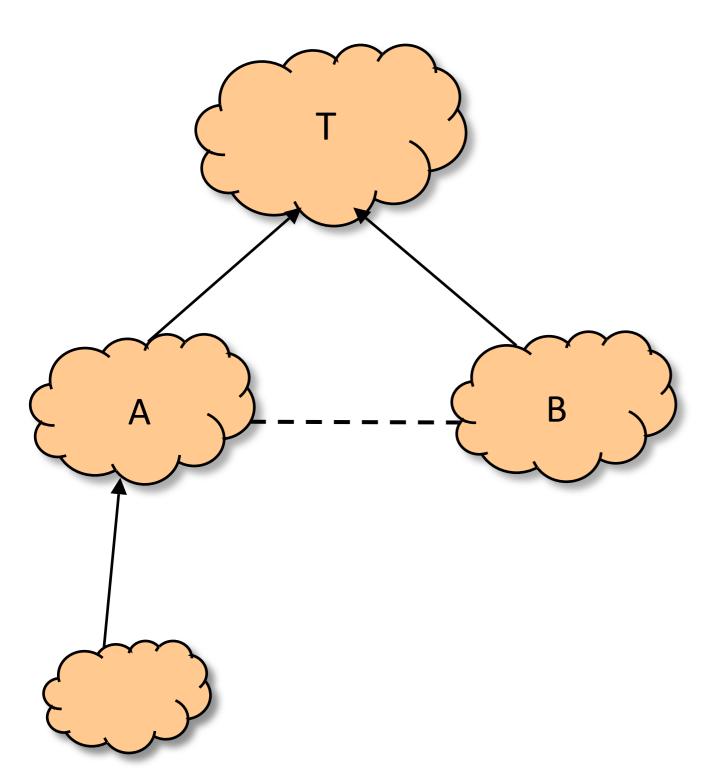




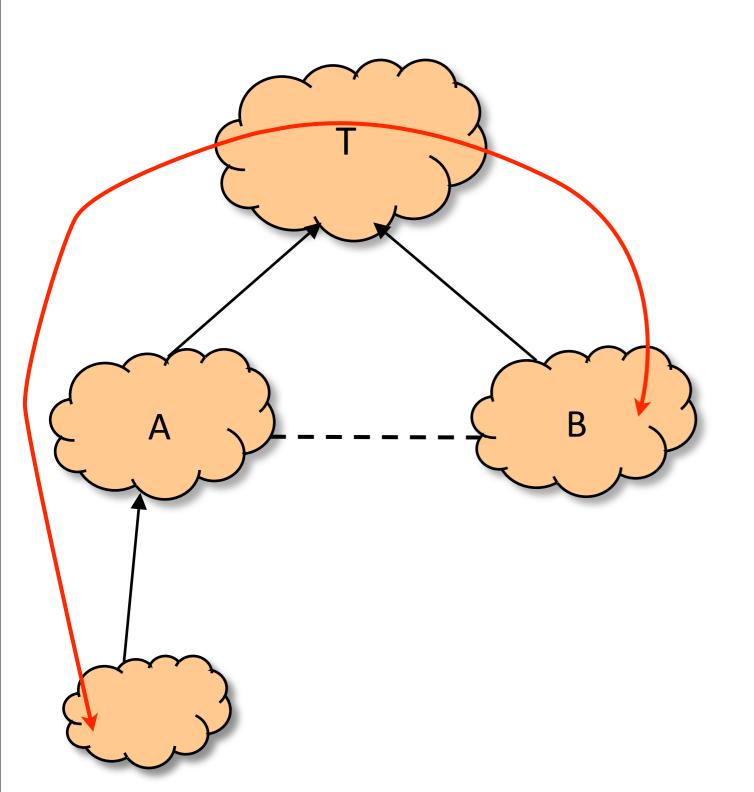
Interconnect cost changes:
 Avoid a transit provider



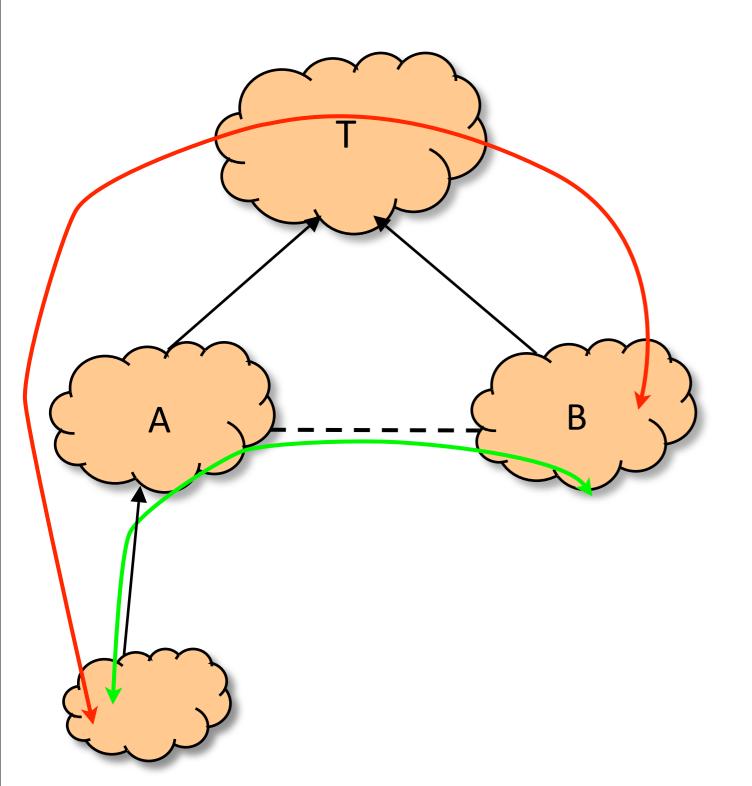
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- Backhaul cost changes:
 Peering link changes how traffic is routed in a network



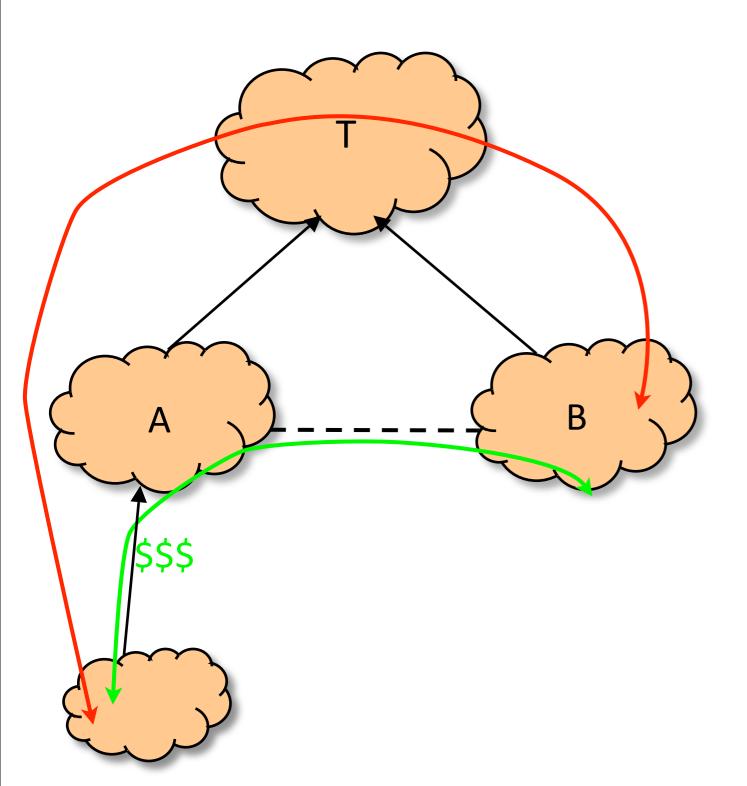
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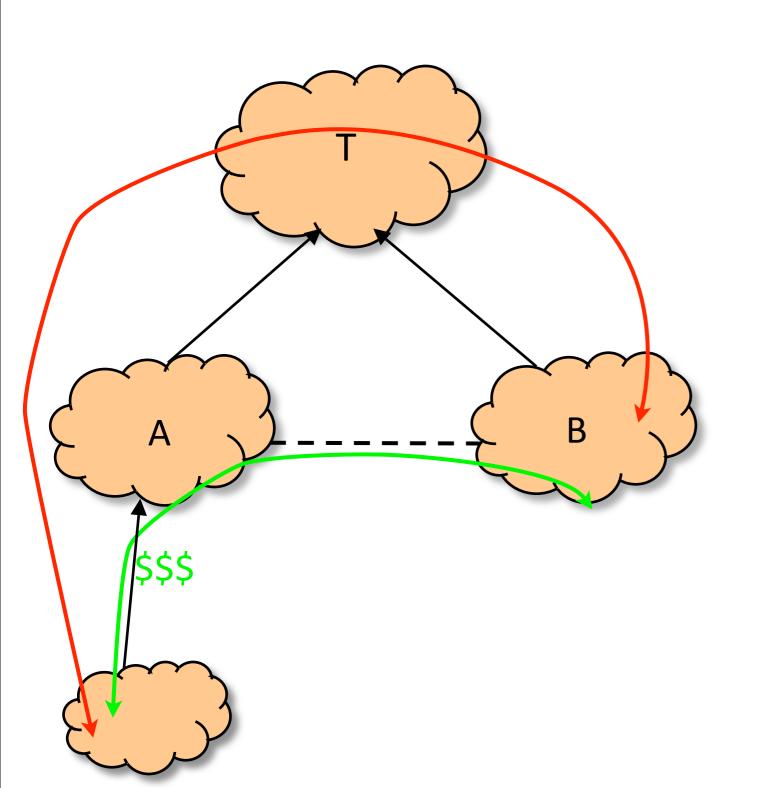
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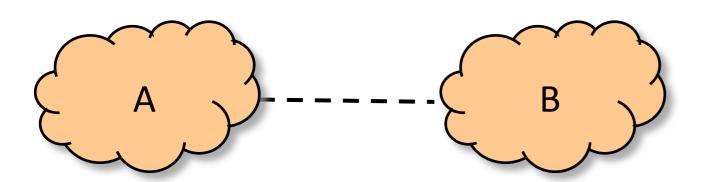
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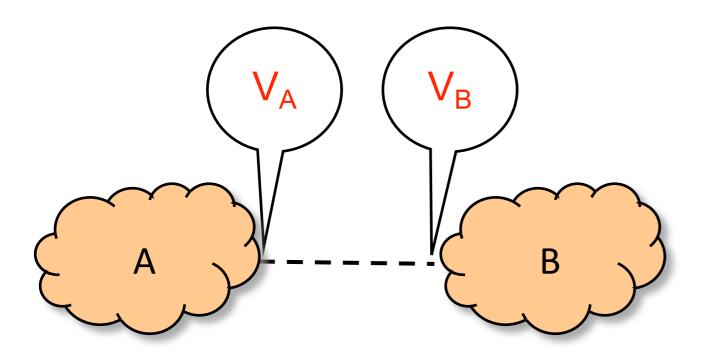


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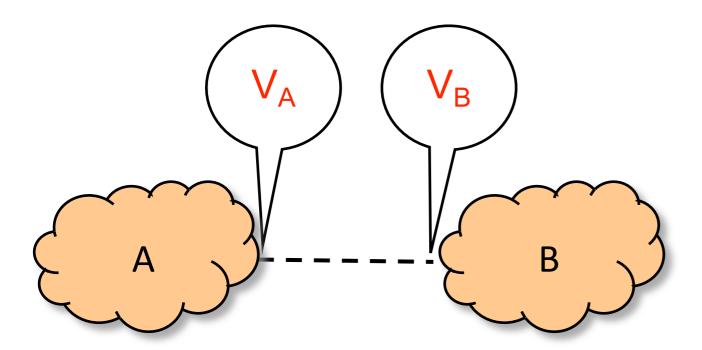


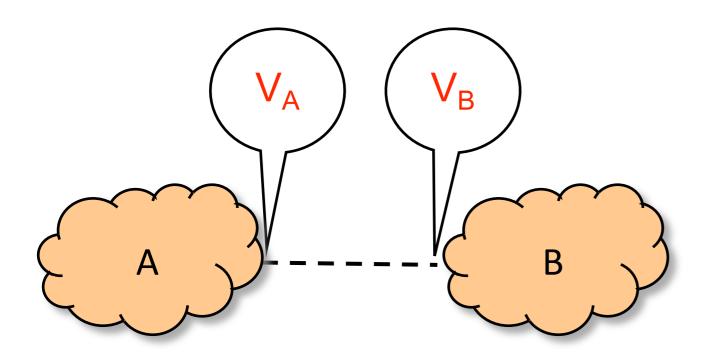
- Interconnect cost changes:
 Avoid a transit provider
- Backhaul cost changes:
 Peering link changes how traffic is routed in a network
- Revenue changes: Attract/ lose traffic due to new peering link





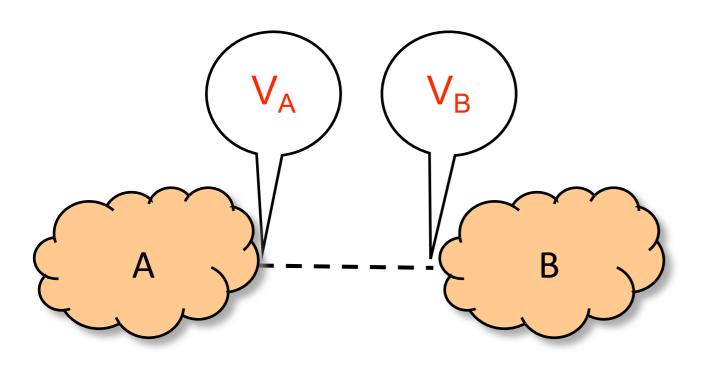
A and B see values V_A and V_B





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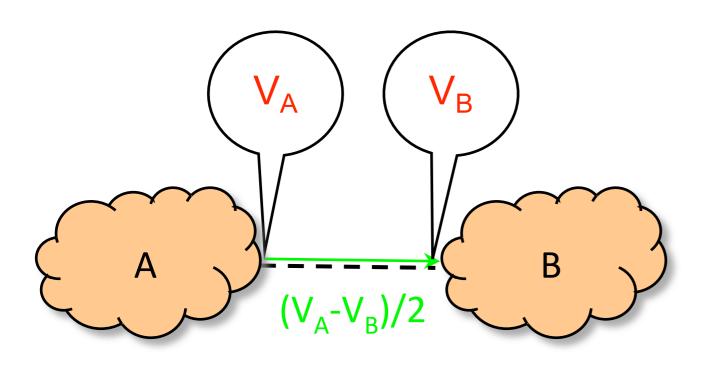
 What should be the paidpeering price?



A and B see values V_A and V_B

 What should be the paidpeering price?

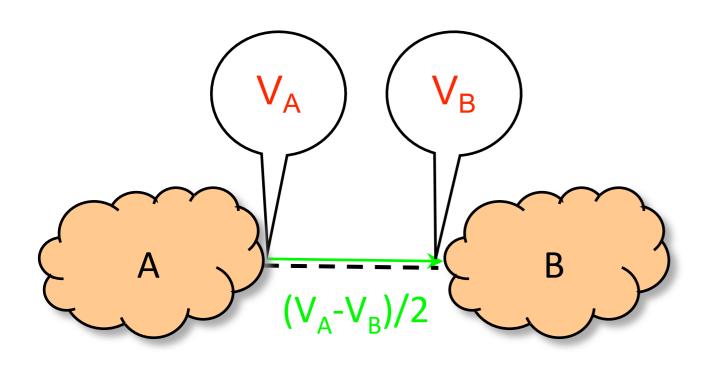
Fair price is (V_A-V_B)/2



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 What should be the paidpeering price?

Fair price is (V_A-V_B)/2

The fair price equalizes the benefit that A and B see from the link

Why Peer at the Fair Price?

- Peering with the fair price is optimal
 - Both networks see better fitness by peering at the fair price
 - As compared to the case where peering link does not exist
- Peering with the fair price is stable
 - No network has the incentive to unilaterally de-peer the other
 - Unique Nash Equilibrium
- Optimal and stable as long as $V_A + V_B > 0$
 - Either V_A or V_B can be negative, as long as total is positive
 - For cost-benefit peering, both V_A and V_B must be positive

Some Hard Questions..

- Value-based peering is fair, optimal and stable.
 - But what (if any) is the right notion of fairness?
 - Equal value? Equal cost?
- How can networks estimate peering value?
 - Peering trials..
 - A cost model to estimate peering value
- What if networks lie about peering value?
- What happens if everyone uses value-based peering?
 - Density of peering links, end-to-end path lengths...

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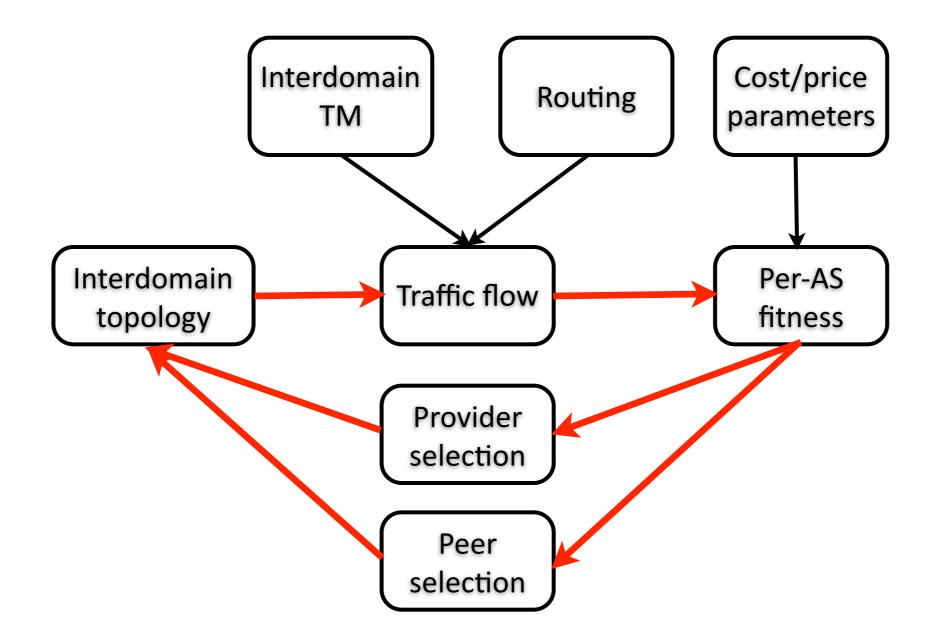
The ITER Model

 ITER: Agent-based computational model to answer "whatif" questions about Internet evolution

• Inputs:

- Network types: transit provider, content provider, stub
- Peer selection methods, provider selection methods
- Geographical constraints
- Pricing/cost parameters
- Interdomain traffic matrix
- Output: Equilibrium internetwork topology, traffic flow, pernetwork fitness

The ITER Approach



 Compute equilibrium: no network has the incentive to change its providers/peers

Using ITER to simulate value-based peering

- Small but realistic internetwork topology: transit providers, content providers, and stubs
- Interdomain traffic matrix: most traffic flows from content providers to stubs
- Provider selection: price-based
- Peer selection: value-based, cost-benefit and traffic-ratio peering

ITER Results: Value-based Peering

- Higher density of peering links with value-based peering as compared to peering by traffic ratios or peering by costbenefit analysis
 - Peering links that cannot be formed with cost-benefit analysis are feasible with value-based peering
 - Shorter end-to-end paths
- Payment direction: The same network can end up on either side of a paid-peering relationship
 - What happens in practice?

Thanks!

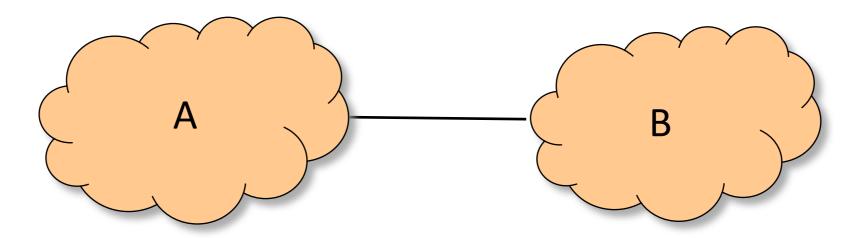
The papers are online at www.caida.org/~amogh Feedback, comments, criticism: amogh@caida.org

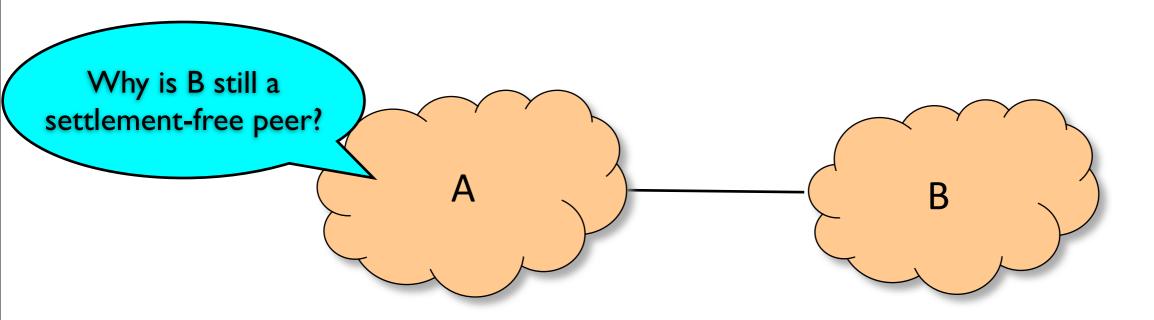
Evaluation

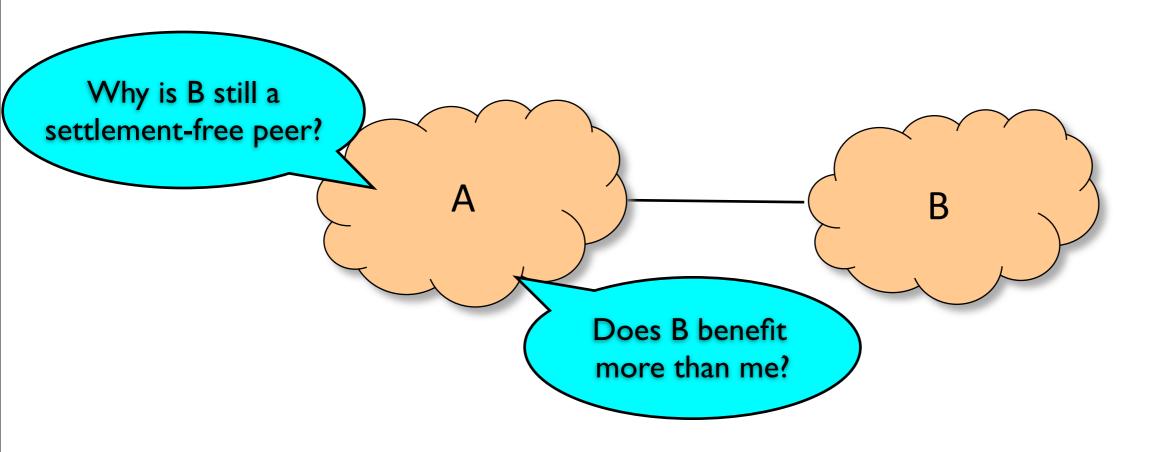
- Access to routing and traffic data from an access ISP in UK
- No access to backhaul and interconnect cost data
- Considered three cost scenarios:
 - Backhaul >> Interconnect (large ISP or cheap transit scenario)
 - Backhaul ≈ Interconnect
 - Backhaul << Interconnect (content provider or expensive transit scenario)
- Evaluated cost optimization using the greedy approach
- What-if scenarios:
 - Where to peer?
 - Which new peer to establish peering with?
 - How useful is an existing peering relation?

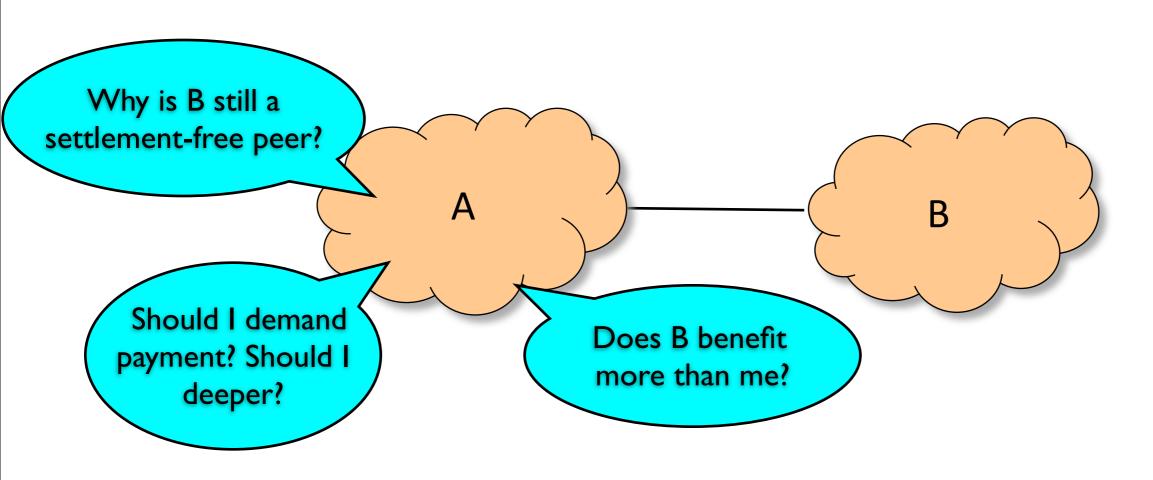
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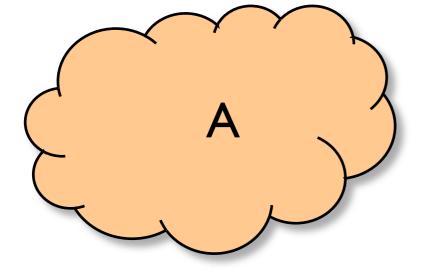




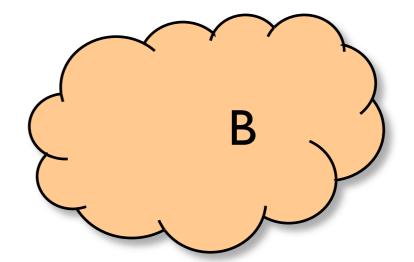


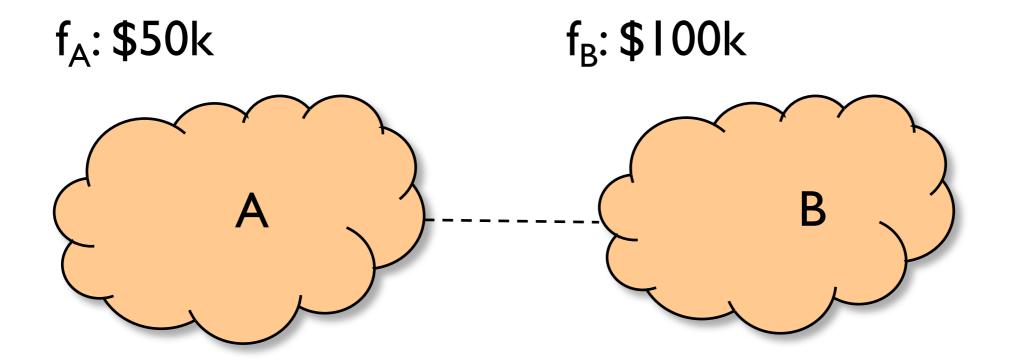


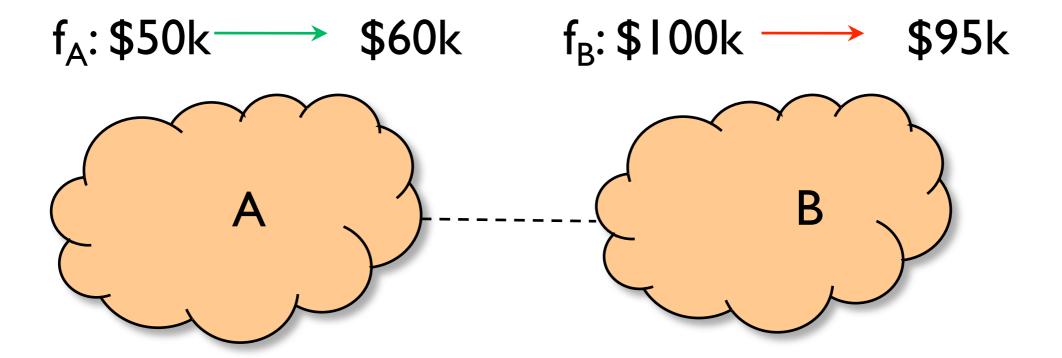
f_A: \$50k

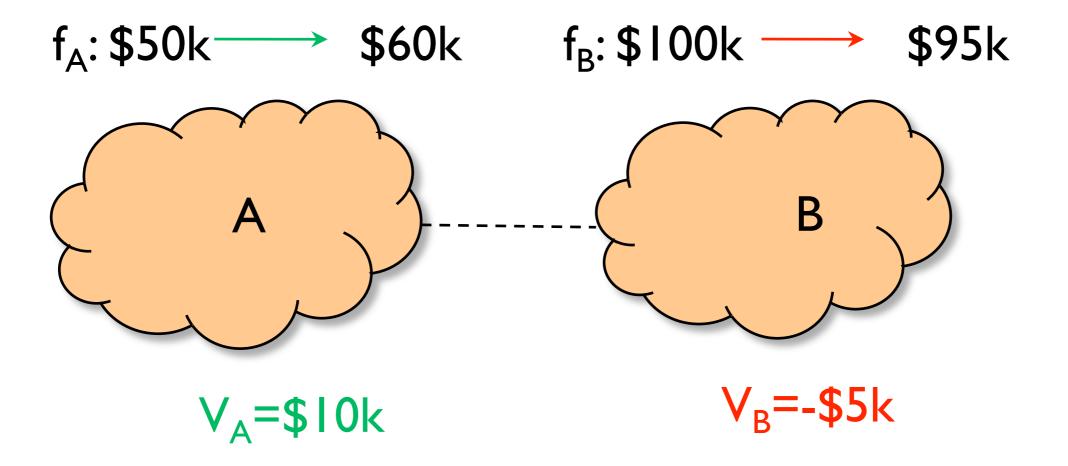


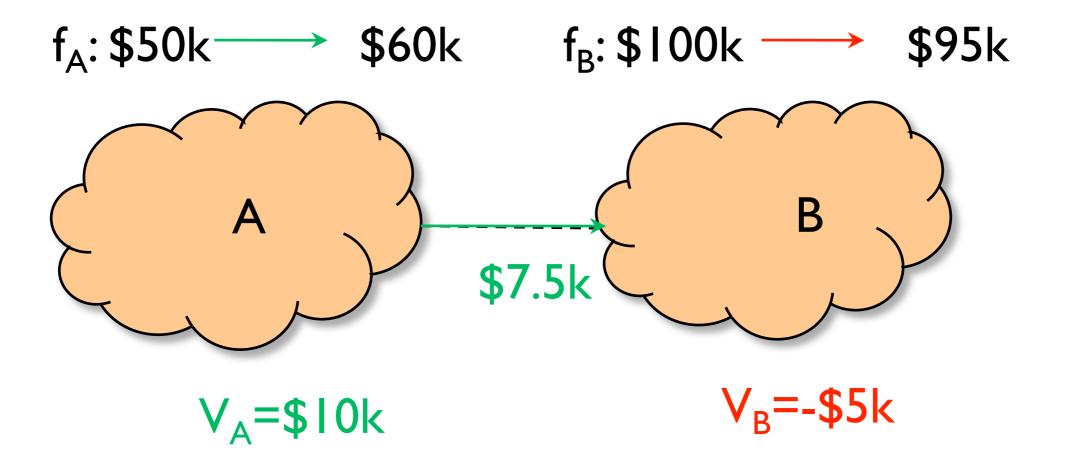
 $f_B: $100k$



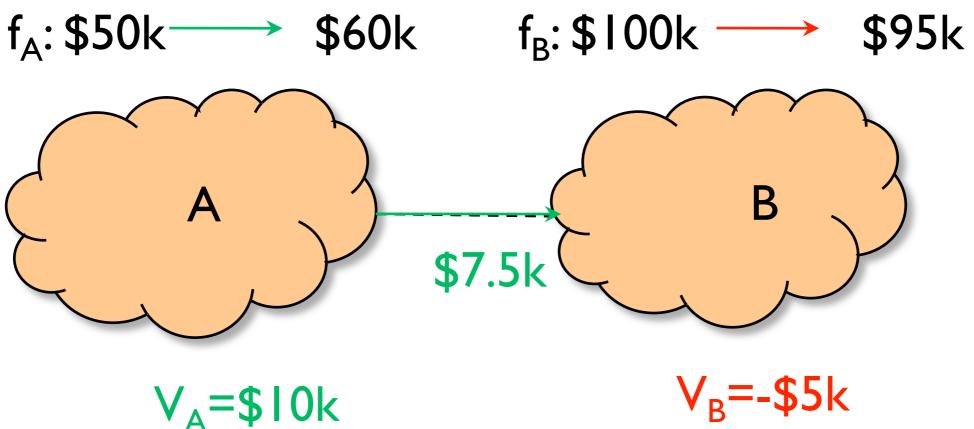


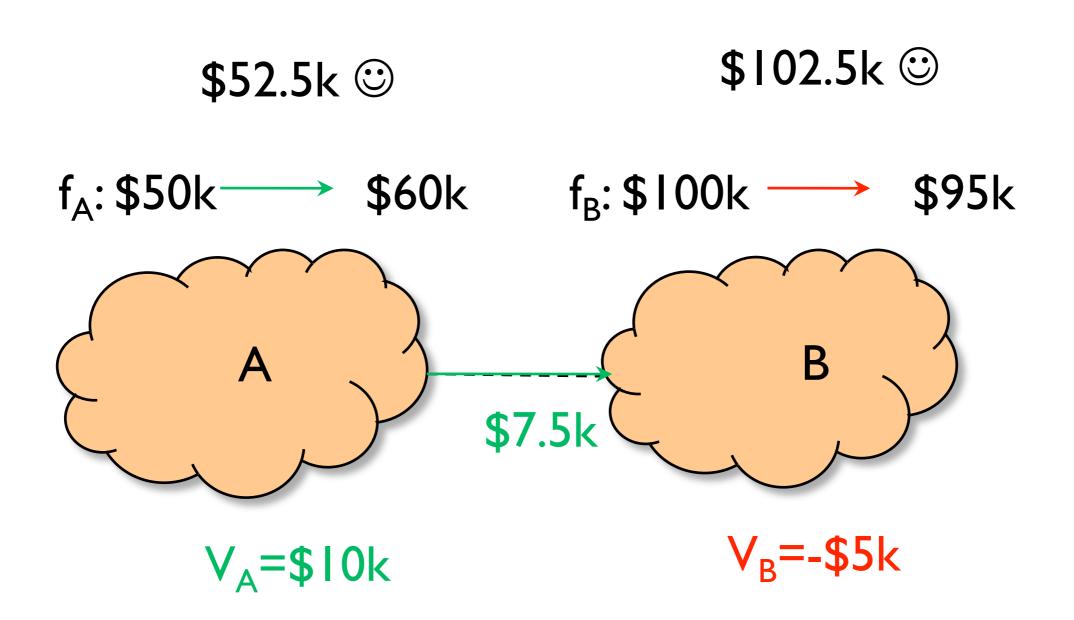






\$52.5k ©





Hiding peering value

- Assume true $V_A + V_B > 0$ and $V_B > V_A$
 - A should get paid (V_B V_A)/2

- If A estimates V_B correctly, and claims its peering value is V_L , where $V_L << V_A$
 - B is willing to pay more: $(V_R V_1)/2$:)

- If A doesn't estimate V_B correctly, and $V_L + V_B < 0$, the peering link is not feasible!
 - A loses out on any payment :(